

# **An Implicit Multidimensional Parity based Data Coding and Decoding Algorithms with FPGA Implementation in Communication Systems**

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**Abstract** –The usability and use of devices in wireless communication is increasing gradually. The mobility and security of the system and the data integrity in the communication process are the requirements of the communication technology for the system. This is the main reason for the development of low voltage generators. Wireless communication medium consists of multiple signals, low voltage communication that can affect any system. Therefore, in wireless systems operating with low-power signals, it is best to use good data encoding and decoding strategies to send data over the medium. The data decoder can recognize changes in the data transmitted in the communication method and correct the data. This encoding-decoding process is called Detection and Correction of error (EDAC) coding. One of such techniques, most often used in communication, is based on the concept of parity encoding and decoding. In the presented study, a multidimensional parity check based coding-decoding algorithm is proposed based on its results and simulation. With the help of proposed work, the associated process can be effectively used to detect and correct 3-bit errors in data transmission.

**Keywords** -Encoder, Decoder, FPGA, HVD, Parity, EDAC, Xilinx.

## **I. INTRODUCTION**

In recent communication, wireless technology has emerged as part of the system. With the development of consumer products, data transmission is also increasing. The transfer of user data is also increasing due to free user-based applications. This provides greater flexibility for mobility and system upgrades than previous technology. Due to the many complex issues in the communication environment, wireless systems are more likely to be affected by noisy signals received by the antennas of the receiving system. Therefore, systems that handle the use of data acquisition signals must have fault detection procedures. Better methods than error detection methods also make corrections to the received data when errors are detected in the received data. Modification of received data can be controlled at the receiving end using the received data in addition to the original data from the sender before sending. However, where the perturbation of the input data is random in nature, it is difficult to choose the best estimate of the actual data. Many methods based on different algorithms have been proposed to solve this error. The main cause of the error is electromagnetic current in the communication channel or in the environment. This signal has the same properties, i.e.to. Noise caused by

signal interference in real signal can be estimated. The amount of noise can be expressed as the number of transmission signals replaced by noise. Many researchers also refer to a small change as a single event upset (SEU) and a large change as a multiple bit upset (MBU). Since looking for interference in data is not only a good solution to problems in the erroneous systems, the system has also been developed to show correction of certain data affecting the bunch of bits. The implementation of this approach can be done at both the hardware and software level. While software systems are used to realize circuits in systems with limited cost, hardware-based circuits are preferred in systems that require higher performance. A simulated method in planning work involving error detection and error correction using four methods. The data is organized in the form of a matrix by creating parity bit streams in following the four matrix bit arrangements, such that - (i) horizontal, (ii) vertical, (iii) forward diagonal and (iv) backward diagonal arrangements. The generated parity bit-stream is written to the data frame at the appropriate time for transmission. The receiver circuits reproduce the parity bits, compare the received parity bits with the generated parity bits to check for errors, and use error correction hardware to complete its job.

## II- LITERATURE REVIEW

The cheapest hardware is always preferred in communication devices with error detection and correction hardware. Few applications of detection of error and correction of error methods are multidimensional (H-V-D horizontal-vertical-diagonal) algorithms. The advantage of this algorithm is that bugs and fixes can be used effectively at both the hardware and software level. Previously, a simple technique to detect errors in a data set was to generate and check for parity. The parity bit can be an odd bit or even slightly different. A bit more of a bit balance only notices bit changes in the bit stream. With single-bit parity, any change in the data stream appears as correct data. This is the impetus for developing methods that use different bits or switch to other data encoding algorithms. An implementation of the multidimensional parity check based data handling scheme with multi-bit error detection and correction scheme in [1], [2], [3], [4] and [5]. A transient error-based detection and correction circuit that reduces the number of inputs and outputs in the connected circuits is shown in [6 and 7]. A similar study is presented in [8] for satellite applications. In [9, 10 and 11], a software-decision-based approach using a

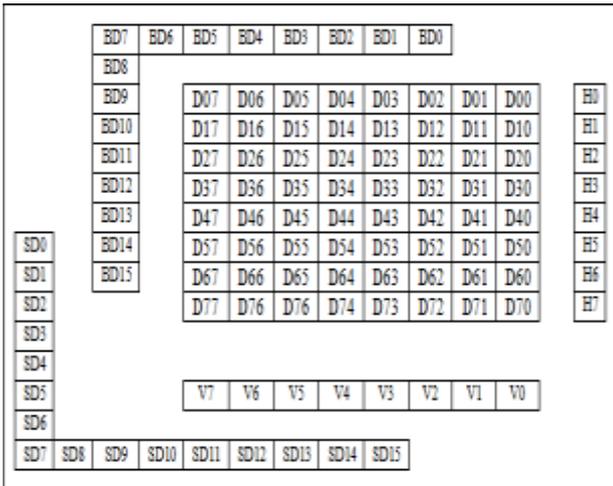
4 D parity based coding scheme is proposed. A 3 bit error detection and correction method is proposed in [12]. Flexible decision-based error correction codes for NAND flash are presented in [13]. H-V-D based fault protection is introduced in [14]. An improved neighborhood error detection system for matrix-based codes is presented in [15]. Cache and memory error detection, correction and mitigation techniques for ground servers and workstations are presented in [16].

## III- METHODOLOGY

The omnidirectional parity code proposed by Methodology is a coding method that has relationships in four directions, such as horizontal, vertical, forward, and backward diagonal. The parity design of the four directions is shown in Figure 1. In this way, the data is first arranged in matrix form. Matrix arrangement of long-length data is described in a particular way so that this arrangement can recognize four different direction based arrangement in the data block. Table 1 shows the data items in the block configuration of the matrix 8X8. Table 1 shows the parity items corresponding to the four different directions. H-V-D parity items are shown in Table 2. The proposed scheme has parity of 8 bits in the horizontal direction, 8 bits in the vertical direction, 15 bits in the diagonal in forward direction and 15 bits in the diagonal in backward direction. For a matrix of size "m x n", the total product in four directions is:

$[m + n + (m + n - 1) + (m + n - 1)] = 2[2(m + n) - 1]$ .  
The encoder input is hardware clock synchronized with the master reset input. The encoder functions to generate 4-odd parity data at the low logic reset control input. If the reset input is logic high, the encoder sets all internal registers and encoder outputs to logic low. The encoder has four equal parts connected together to produce the same output. Each parity generator block is dedicated to generating parity for a particular direction. The functional diagram of the proposed encoder is shown in figure 2(a). The decoder input is a clock synchronous hardware design that is also a master reset control input. The operating diagram of the decision maker is shown in figure 2(b). The decoder operates on the received data and the equivalent bit of logic low resets the control input. If the reset input is logic high, the decoder sets all internal registers and the decoder's output to a logic without value. When enabled for processing, the decoder first generates horizontal, vertical, forward and backward diagonal parity bits from the received data bits. The decoder then compares the received parity bits with the generated parity bits. If the two parity sets are found to be unequal, the output error

is set to a high value, otherwise a low value. If the error correction control input is enabled, the data is sent to the error correction logic block. If error correction causes the input to be set to logic low, no error correction causes data to be received. In this case, only information about the presence or absence of an error will be sent to the output of the previous block using the error indicator signal.



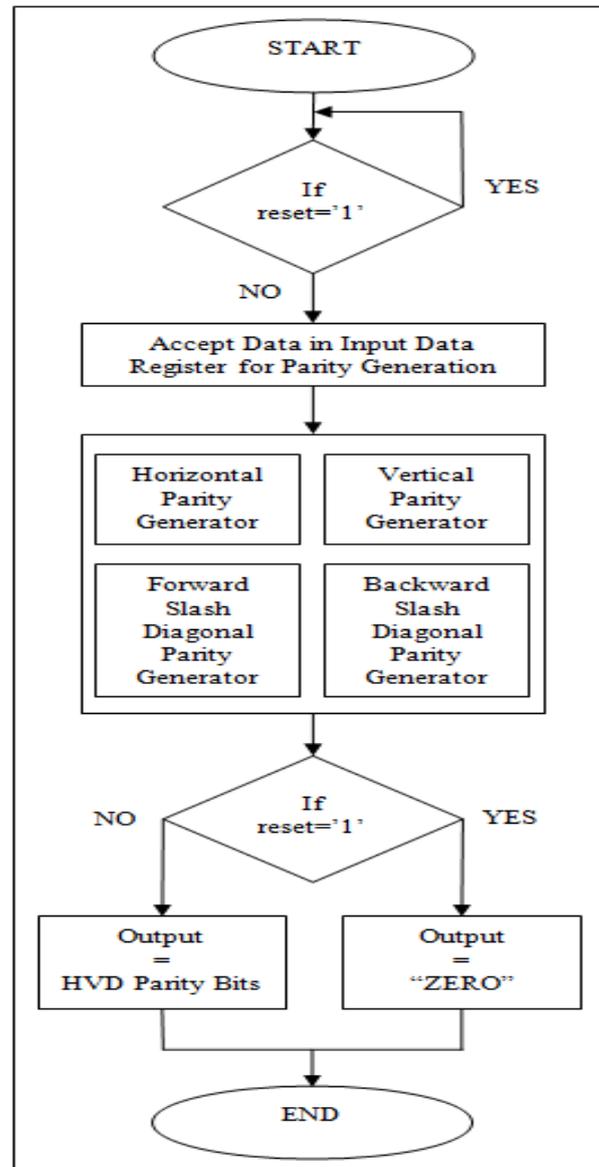
**Fig. 1:** Horizontal-Vertical-Diagonal H-V-D Parity based check method in 8X8 matrix form

**Table 1:** Data blocks in 8X8 matrix form arranged in rows

Row Number	Arranged Data Bits from MSB to LSB
1	D07-D06-D05-D04-D03-D02-D01-D00
2	D17-D16-D15-D14-D13-D12-D11-D10
3	D27-D26-D25-D24-D23-D22-D21-D20
4	D37-D36-D35-D34-D33-D32-D31-D30
5	D47-D46-D45-D44-D43-D42-D41-D40
6	D57-D56-D55-D54-D53-D52-D51-D50
7	D67-D66-D65-D64-D63-D62-D61-D60
8	D77-D76-D75-D74-D73-D72-D71-D70

**Table 2:** Generation of Parity Bits using Encoder

Generation of parity bits using Encoder	
<b>Horizontal Parity Bits</b>	H0, H1, H2, H3, H4, H5, H6, H7
<b>Vertical Parity Bits</b>	V0, V1, V2, V3, V4, V5, V6, V7
<b>Forward Diagonal Parity Bits</b>	SD0, SD1, SD2, SD3, SD4, SD5, SD6, SD7, SD8, SD9, SD10, SD11, SD12, SD13, SD14, SD15
<b>Backward Diagonal Parity Bits</b>	BD0, BD1, BD2, BD3, BD4, BD5, BD6, BD7, BD8, BD9, BD10, BD11, BD12, BD13, BD14, BD15



**Fig. 2:** (a) Encoder Block wise operation

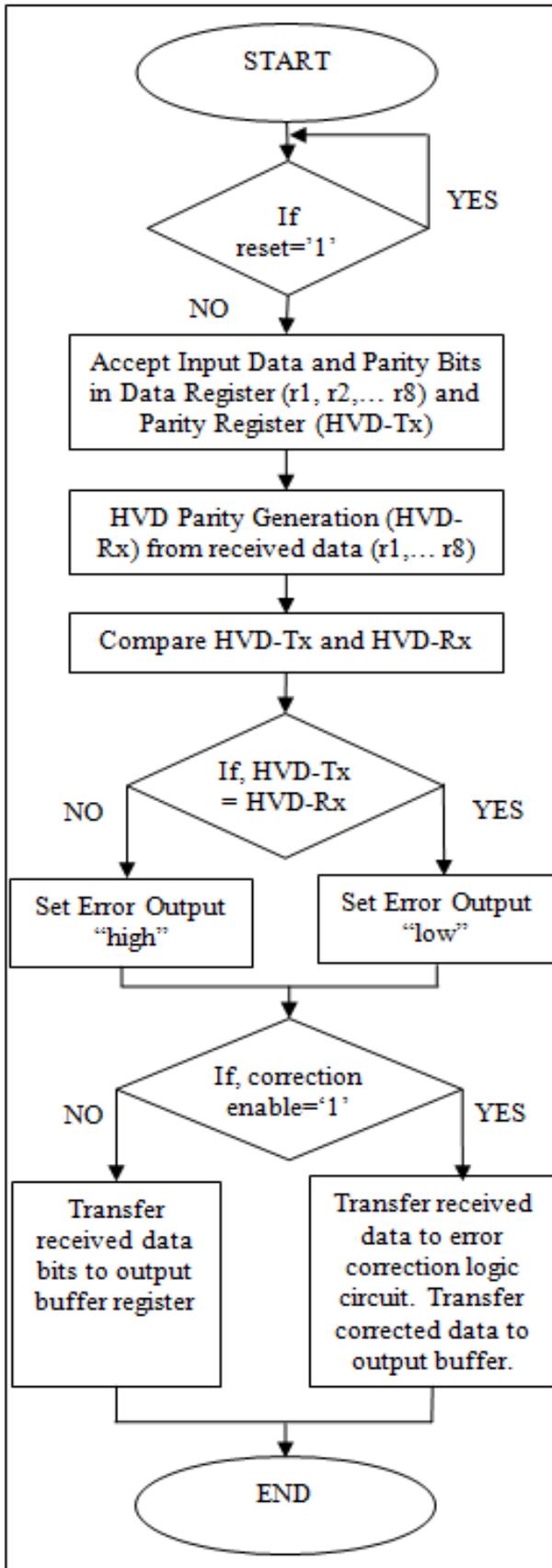


Fig. 2:(b) Decoder Block wise operation

IV- SYNTHESIS&SIMULATIONOUTCOMES

The simulation of proposed H-V-Dtest bench design is done using Xilinx ISE Toolusing VHDL Language environment. The result showing waveforms of test bench simulation of encoder part is as shown in figure 3. The generated output parity bits along with binary data streams of input of receiver side are as follows –

- Row No. 1 = 10011011
- Row No. 2 = 01100110
- Row No. 3 = 10101010
- Row No. 4 = 11000111
- Row No. 5 = 10000001
- Row No. 6 = 01011010
- Row No. 7 = 01000011
- Row No. 8 = 11100100

- Horizontal Parity Bits = 01001001
- Vertical Parity Bits = 11101100
- Forward Diagonal Parity Bits = 110010101010111
- Backward Diagonal Parity Bits = 010011000001001

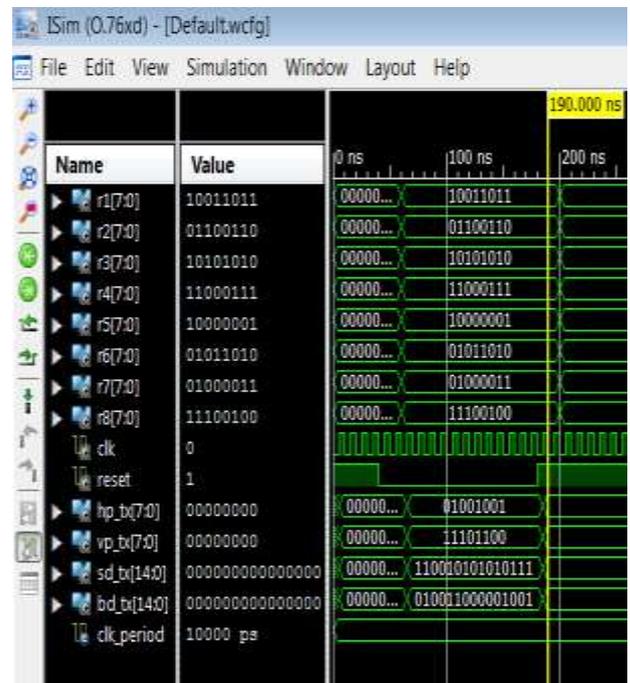


Fig. 3: Simulated Output of Encoder

The output waveforms of simulation of decoder are as shown in figures – 4 & 5. These waveforms are of following two cases:

- Case I: when there is no error in the data received, and
- Case II: when 4-bit error is introduced and correction of

erroneous data in the received data by the decoder. The inputs to decoder for simulation as given above two cases are as shown in following Table – 3:

**Table 3: Simulation Inputs of Proposed Decoder**

Inputs to Decoder	Case I	Case II
Row No. 1	10011011	10011011
Row No. 2	01100110	01100110
Row No. 3	10101010	10101010
Row No. 4	11000111	11000111
Row No. 5	10000001	10000001
Row No. 6	01011010	01011010
Row No. 7	01000011	01 <b>1111</b> 11
Row No. 8	11100100	11100100

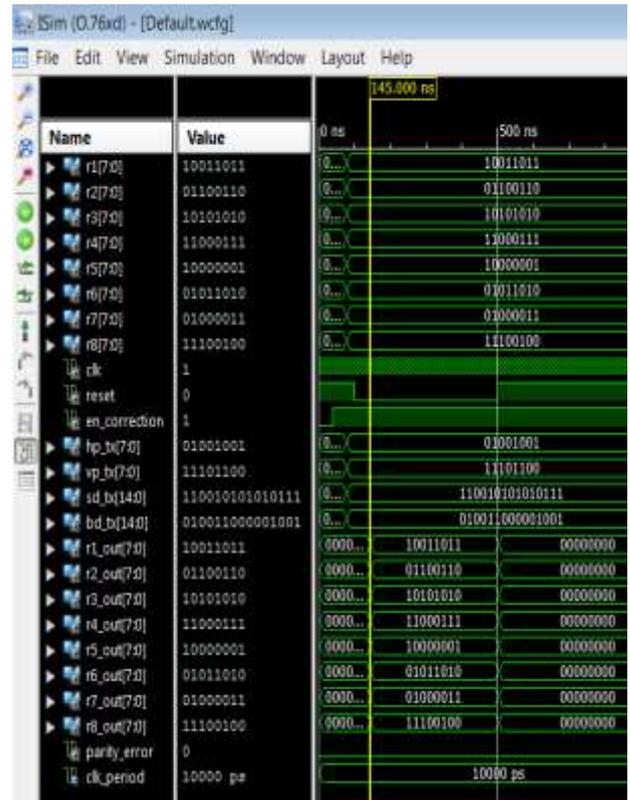
Xilinx Synthesis Tool on Xilinx Spartan 3E XC3S500E-4FG320 FPGA device platform is used for performing synthesis of hardware based design proposed. Following tables 4 & 5 shows the optimum hardware utilization summary for proposed encoder & decoder designs.

**Table 4: Utilization summary of Hardware of proposed Encoder**

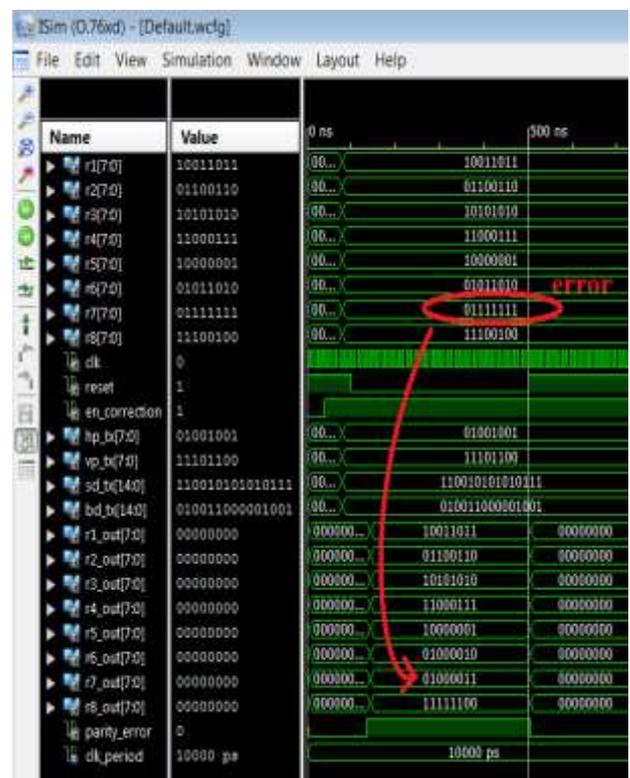
Hardware Resources used	Total	Encoder Utilization	
		Used	%
No. of Slices	4656	47	1
No. of LUTs	9312	68	1
No. of Flip-flops	9312	90	1

**Table 5: Utilization summary of Hardware of proposed Decoder**

Hardware Resources used	Total	Decoder Utilization	
		Used	%
No. of Slices	4656	227	4
No. of LUTs	9312	184	1
No. of Flip-flops	9312	329	3



**Fig. 4: Simulated output for Case I of Decoder**



**Fig. 5: Simulated output for Case II of Decoder**

## V-RESULT & DISCUSSION

In this method of work, a technique for error detection and correction is proposed. The system design examines and corrects all sequential and multiple errors from data blocks less than 4 bits in length up to 3 bits. In exceptional cases the 4 bit error is also within the repair capacity of this algorithm. The execution strategy also generates output signals with error signals showing indications of errors. This output stimulates the interface circuitry to take decision whether to accept or reject the data received.

## VI- CONCLUSION

Omnidirectional parity check algorithm plays an important role in detection and correction of errors in data communication, so the proposed work using the FPGA algorithm shows that the data capacity is improved and optimum hardware resources are utilized. The proposed algorithm can be improved in the future to correct all 3 bit and up to 4 bit errors. Additionally, the proposed operation can be improved by running simulations to detect and correct burst errors. This working concept can also be combined with other techniques to identify errors and correct the process as an attempt to improve the system.

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